Integer Programming Solvers

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August 16, 2008

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 - Integer Programming (IP)
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Motivation

- Combinatorial optimisation, combinatorial explosion
- Universal "solver"
- Solvers sacrificing optimality (local search)
- Solvers sacrificing performance (?!)
- Solvers and their "workhorses"

Linear Programming

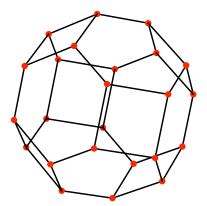
The Problem of Linear Programming

objective : max
$$c_1x_1+c_2x_2+\cdots+$$
 c_nx_n constraints : $a_{1,1}x_1+a_{1,2}x_2+\cdots+$ $a_{1,n}x_n\leq b_1$ $a_{m,1}x_1+a_{m,2}x_2+\cdots+$ $a_{m,n}x_n\leq b_m$ $x_1,\ldots,x_n\geq 0$

- Modern solvers (CPLEX) run almost in linear time on average
- Structured instances with hundreds of millions of variables can be solved in hours

Linear Programming (Cntd.)

- Simplex: hopping over the corner points
- Interior point methods: go straight
- Ellipsoid methods: of theoretical interest



Integer Programming (IP)

The Problem of Integer Programming

objective : max
$$c_1x_1+c_2x_2+\cdots+$$
 c_nx_n constraints : $a_{1,1}x_1+a_{1,2}x_2+\cdots+$ $a_{1,n}x_n\leq b_1$ $a_{m,1}x_1+a_{m,2}x_2+\cdots+$ $a_{m,n}x_n\leq b_m$ $x_1,\ldots,x_n\in \mathbb{Z}^+$

- Where is the difference?
- How important it is?

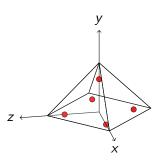


Figure: A poor linear programming relaxation.

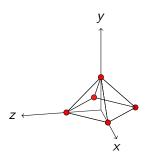


Figure: The convex hull of integer points. Typically, there are exponentially many facets.

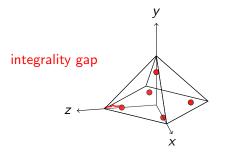


Figure: The "integrality gap" is the distance from the LP optimum to the IP optimum – or the value of the present best integer feasible solution.

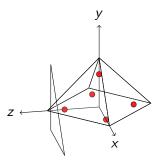


Figure: Pure cutting plane methods can go for ever ...

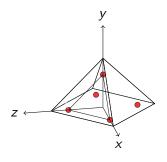


Figure: ... even if your cuts are facet-defining. Why?

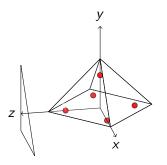


Figure: Branching can be thought of as a form of "divide and conquer". Here, branch on a variable by adding x > 2 ...

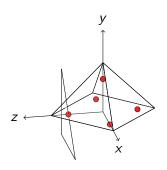


Figure: ... or by adding $x \le 1$.

- Branch and bound: Building, pruning and traversing a tree
- Cuts by non-trivial rounding or polytope specific
- Loads of heuristics
- Branching heuristics: what (variable) to branch on, which direction to try first
- Primal heuristics: find the first feasible solution so that you can prune the tree
- Improvement heuristics: very often the source of feasible solutions
- Search strategies: which node in the tree should be visited next
- Cut generation: how much time to spend in separation, what cuts are good enough, when should they be removed, ...

What I am doing at Nottingham

Formerly:

- Preprocessing of a discrete set for polylog extreme ray queries
- Optimal transformation of graph colouring to multicolouring and applications

More recently:

- Primal heuristics for structured instances of binary integer programming
- Branching heuristics for binary integer programming
- Auto-tuning integer programming solvers

What I Was Doing Here

- An applet using VTK Java Wrappers
 Sends JVM crashing down
- Snippet for SVG and PDF generation from Jython Turns out to be rather bulky
- Suggestion: Sketch/TikZ output from plot3d Nils implemented a preliminary version Sketch turns out to be rather lame

http://wiki.sagemath.org/JakubMarecek

Conclusions

- For inapproximable problems, you need "heuristics"
- Design of "heuristics" often requires some insight
- Theoretical analysis is difficult, visualisations would help

See also:

- Sven Krumke: Integer Programming optimierung.mathematik.uni-kl.de/~krumke/
- Lex Schrijver: A Course in Combinatorial Optimization
 www.cwi.nl/~lex/
- Vašek Chvátal: Linear Programming
- Nemhauser & Wolsey: Integer and Combinatorial Opt.
- H. P. Williams: Model Building in Mathematical Programming

Many thanks!

- Many thanks to the organisers and you all!
- Any questions or comments are appreciated!
- My email addresses are at http://cs.nott.ac.uk/~jxm