

**Some ideas for efficient implementation of
algorithms for polynomial matrix computations**

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Introduction: integers versus univariate polynomials

Integers:

- well suited ring for a binary computer
 - array of limbs of 64 bits each
- primes are quite dense
 - prime number theorem: there are about $N/\ln N$ primes $\leq N$
 - more than 10^{15} 64-bit primes
 - enough for all practical purposes

Polynomials $K[x]$:

- dense, sparse, supersparse (lacunary)?
- what is the coefficient field K ?
 - $K = \mathbb{Z}/(p)$, $K = \text{GF}(2)$, $K = \mathbb{Q}$
 - K an extension field of one of the above
- and then... what about multivariate? coefficient ring not a field?

The core problem: Polynomial matrix multiplication

Representation:

- matrices of polynomials:

$$F = \begin{bmatrix} x+1 & 4x^2+3x \\ 4x^2+3x+1 & 4x^2+6x+2 \end{bmatrix}$$

- polynomials with matrix coefficients:

$$F = \begin{bmatrix} 0 & 4 \\ 4 & 4 \end{bmatrix} x^2 + \begin{bmatrix} 1 & 3 \\ 3 & 6 \end{bmatrix} x + \begin{bmatrix} 1 & 0 \\ 1 & 2 \end{bmatrix}$$

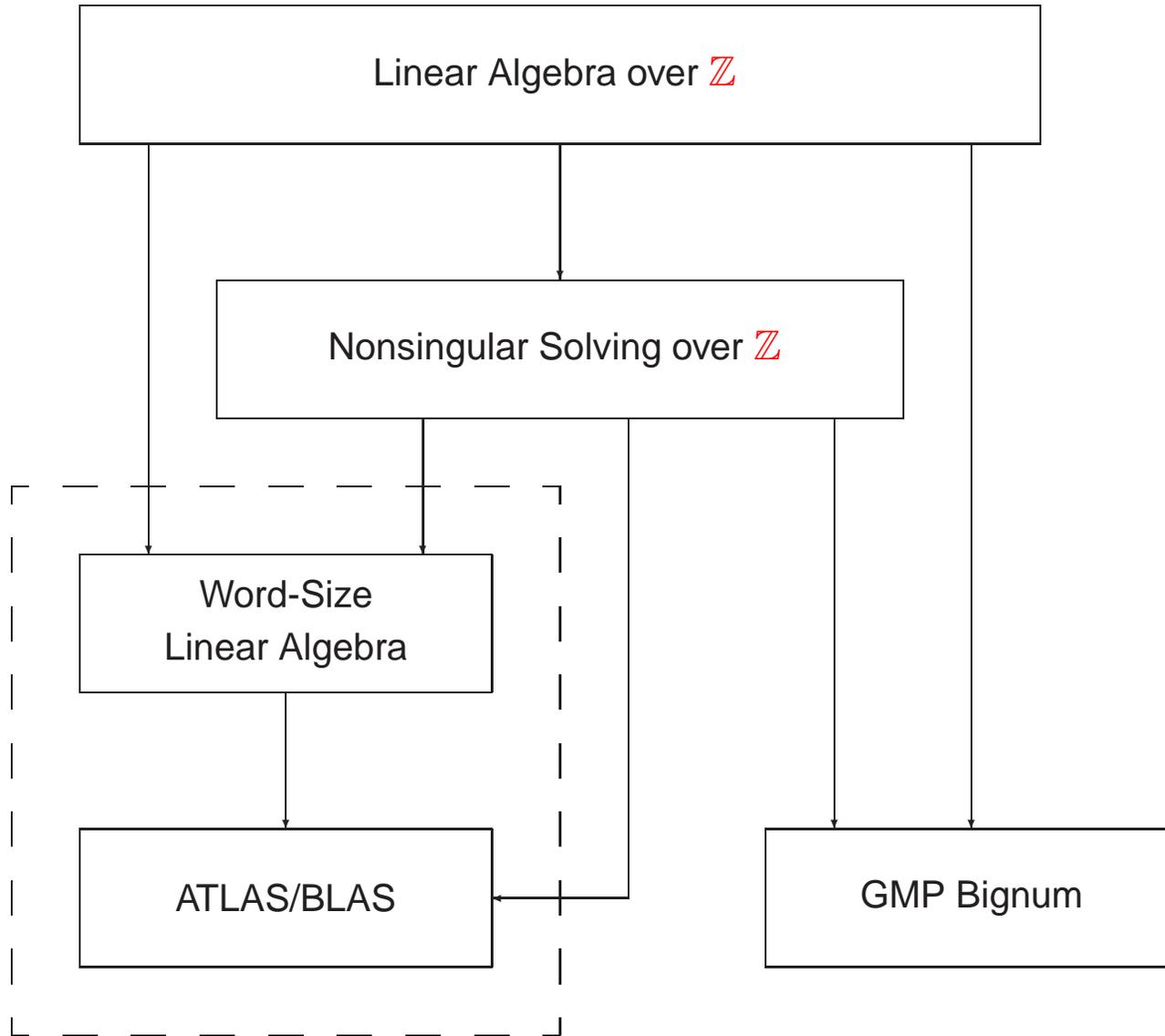
- compact:

$$F = \left[\begin{array}{cc|cc|cc} 0 & 4 & 1 & 3 & 1 & 0 \\ 4 & 4 & 3 & 6 & 1 & 2 \end{array} \right]$$

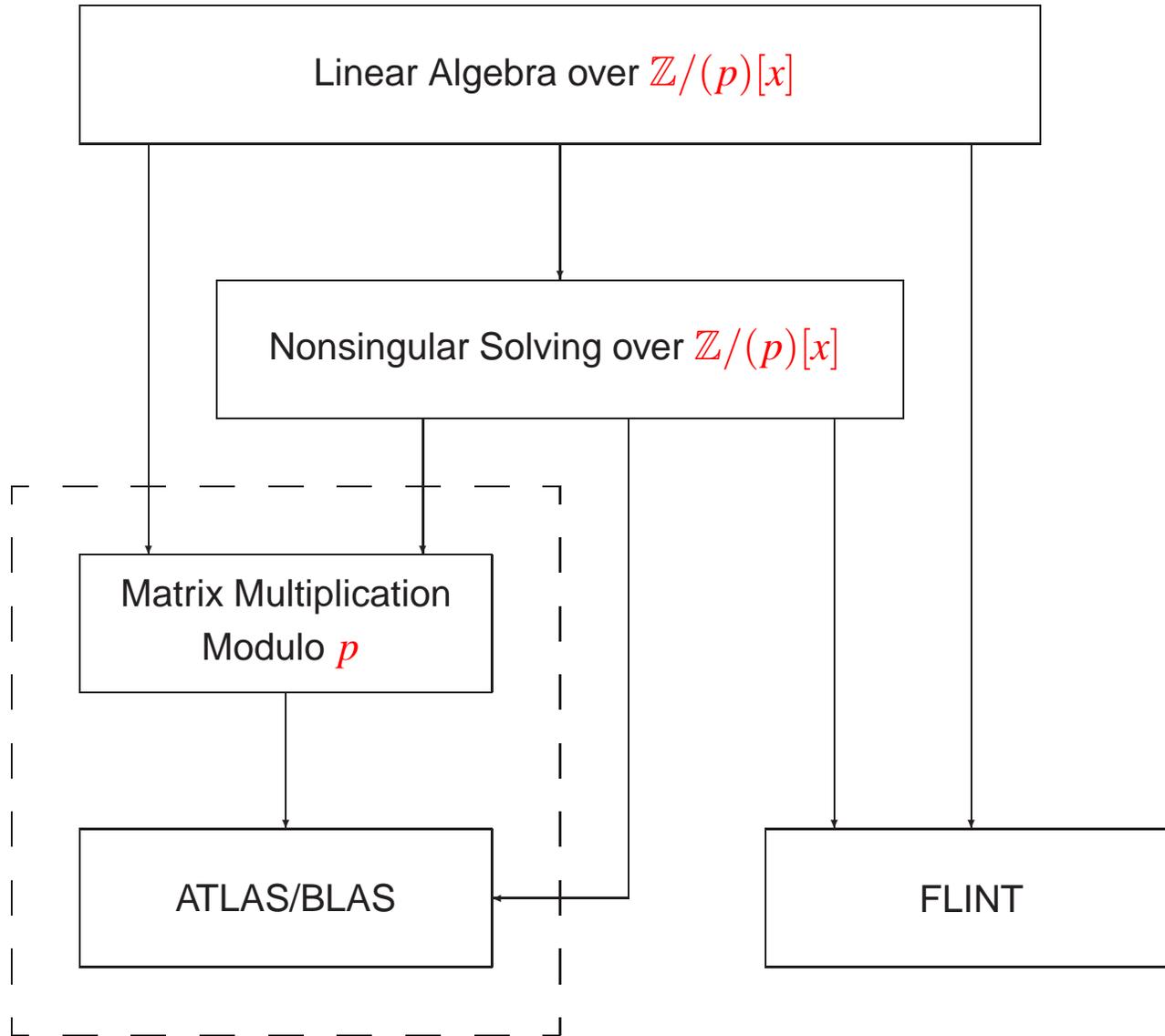
Methods:

1. triply nested for loop using polynomial multiplication (easiest)
2. via integer matrix multiplication: bit packing (FLINT [Frederik J.])
3. evaluation and interpolation (asymptotically fastest)
 - in-place truncated FFT [Harvey and Roche, 2010]?

Organization of the Integer Matrix Library: IML



Organization of the Polynomial Matrix Library: PML?



Example of linear solving: $K = \mathbb{Z}/(7)$, $n = 5$, $d = 2$

Input:

$$A = \begin{bmatrix} 6x^2 + 6x + 4 & 2x^2 + 2x + 5 & 3x^2 + 2x + 1 & 5x + 3 & 2x^2 + 6x + 6 \\ 4x^2 + 4x & 2x^2 + 5x + 3 & 5x + 4 & 3x^2 + 5x & 4x^2 + x + 1 \\ 6x^2 + 2 & 2x^2 + x + 6 & 4x^2 + 2x + 2 & 4x^2 + 5x + 1 & x^2 + x \\ 3x^2 + x + 4 & 5x + 6 & 4x^2 + 2x + 1 & 2x^2 + 6x + 2 & 3x^2 + x \\ x^2 + 2x + 6 & 2x^2 + 5x + 5 & 4x^2 + 4x & 6 & x^2 + x + 3 \end{bmatrix} \quad b = \begin{bmatrix} x^2 + x + 2 \\ 5x^2 + x + 2 \\ 3x^2 + 3x + 5 \\ x^2 + 5x + 3 \\ 2x + 2 \end{bmatrix}$$

Output:

$$v := A^{-1}b = \begin{bmatrix} \frac{4x^9 + 2x^8 + 5x^7 + 6x^6 + x^5 + 3x^4 + 6x^3 + 2x + 2}{4x^{10} + x^9 + 3x^8 + 4x^7 + 4x^6 + 2x^5 + 6x^4 + 6x^3 + 2x^2 + 5x + 3} \\ \frac{6x^{10} + 2x^9 + x^7 + 3x^6 + x^5 + 6x^4 + 3x^3 + 6x^2 + 2x + 1}{4x^{10} + x^9 + 3x^8 + 4x^7 + 4x^6 + 2x^5 + 6x^4 + 6x^3 + 2x^2 + 5x + 3} \\ \frac{x^{10} + 5x^9 + 6x^8 + 3x^6 + 5x^5 + 3x^4 + 5x^3 + x^2 + 3x + 5}{4x^{10} + x^9 + 3x^8 + 4x^7 + 4x^6 + 2x^5 + 6x^4 + 6x^3 + 2x^2 + 5x + 3} \\ \frac{3x^{10} + 3x^9 + 6x^8 + 6x^7 + 5x^6 + 5x^4 + 6x^3 + x^2 + x + 2}{4x^{10} + x^9 + 3x^8 + 4x^7 + 4x^6 + 2x^5 + 6x^4 + 6x^3 + 2x^2 + 5x + 3} \\ \frac{5x^{10} + 4x^9 + 5x^8 + 2x^7 + x^6 + 2x^5 + 5x^4 + 5x^3 + 3x^2 + 6x + 4}{4x^{10} + x^9 + 3x^8 + 4x^7 + 4x^6 + 2x^5 + 6x^4 + 6x^3 + 2x^2 + 5x + 3} \end{bmatrix}$$

Outline of Y -adic lifting for system solving

1. Radix expansion of solution: $Y = x^2$

$$\frac{5x^2 + 6x + 3}{x^2 + 4x + 3} \equiv (1 + 3x) + (2 + x)Y + (1 + 5x)Y^2 \pmod{Y^3}$$

2. Radix conversion:

$$(1 + 3x) + (2 + x)Y + (1 + 5x)Y^2 = 1 + 3x + 2x^2 + x^3 + x^4 + 5x^5$$

3. Rational function reconstruction:

$$\frac{5x^2 + 6x + 3}{x^2 + 4x + 3} \equiv 1 + 3x + 2x^2 + x^3 + x^4 + 5x^5 \pmod{x^5}$$

Nonsingular rational system solving over $K[x]$ via lifting

Input: $A \in K[x]^{n \times n}$ and $b \in K[x]^{n \times 1}$ of degree d

Compute: $A^{-1}b \in K(x)$

Method:

1. Choose $Y \in K[x]$ such that $\gcd(Y, \det A) = 1$

Set $k = \lceil 2nd / \deg Y \rceil$

2. Compute $B = \text{Rem}(A^{-1}, Y)$

3. $r := b$

for $i = 0$ **to** $k - 1$ **do**

$v_i := \text{Rem}(Br, Y)$ # $\deg v_i < \deg Y$

$r := (r - Av_i) / Y$ # $\deg r < d$

4. Reconstruct $A^{-1}b$ from $A^{-1}b \equiv v_0 + v_1Y + \cdots + v_{k-1}Y^{k-1} \pmod{Y^k}$

→ what should degree of Y be?

→ what should factorization of Y be?

Lifting using a “lifting basis”

1. Select of modulus:

$$Y = (x - \alpha_1)(x - \alpha_2) \cdots (x - \alpha_{\deg Y})$$

$$Z = (x - \beta_1)(x - \beta_2) \cdots (x - \beta_d) \text{ with } \gcd(Z, Y) = 1$$

2. Initialization:

$$\text{old: } B := \text{Rem}(A^{-1}, Y)$$

$$\text{new: } (B_1, \dots, B_{\deg Y}) := ((A|_{x=\alpha_1})^{-1}, \dots, (A|_{x=\alpha_{\deg Y}})^{-1})$$

$$(C_1, \dots, C_d) := (A|_{x=\beta_1}, \dots, A|_{x=\beta_d})$$

3(a). Lifting step:

$$\text{old: } v_i := \text{Rem}(Br, Y)$$

$$\text{new: } ((v_i)|_{x=\alpha_1}, \dots, (v_i)|_{x=\alpha_{\deg Y}}) := (B_1 r|_{x=\alpha_1}, \dots, B_{\deg Y} r|_{x=\alpha_{\deg Y}})$$

3(b). Residue update:

$$\text{old: } r := (r - Av_i)/Y$$

$$\text{new: } r|_{x=\beta_j} := (r|_{x=\beta_j} - C_j(v_i)|_{x=\beta_j}) \text{Rem}(Y^{-1}, x - \beta_j) \text{ for } j = 1, \dots, d$$

Lifting using a “lifting basis”: main work

2. Initialization:

→ $\text{deg } Y$ matrix inversions over K

3(a). Lifting steps:

→ total $2nd$ matrix \times vector products over K

3(b). Residue updates:

→ total $2nd \times (d / \text{deg } Y)$ matrix \times vector products over K

4. Reconstruct solution using interpolation + radix conversion + rational function reconstruction

Key optimizations:

- Choose $\text{deg } Y$ to balance costs of phases 2 and 3.
→ can be automatically tuned
 - Reduce the $2nd$ using vector rational function reconstruction.
→ decrease $2nd$ but increase cost of step 4
- [Olesh & Storjohann, 2007]

Linearization for polynomial lattice basis reduction

First consider Euclidean algorithm over $\mathbb{Z}/(7)[x]$

$$\begin{bmatrix} 4x^3 + 6x^2 + 5x + 6 \\ 4x^3 + 2x^2 + 3x + 5 \end{bmatrix} \rightarrow \begin{bmatrix} 4x^3 + 6x^2 + 5x + 6 \\ 3x^2 + 5x + 6 \end{bmatrix} \rightarrow \begin{bmatrix} 4x^2 + 4x + 6 \\ 3x^2 + 5x + 6 \end{bmatrix} \rightarrow \dots \rightarrow \begin{bmatrix} x + 6 \\ 0 \end{bmatrix}$$

Same idea works for lattice reduction of matrices

(Note: $[d] \equiv$ a polynomial of degree d)

$$\begin{array}{c} A \\ \left[\begin{array}{cccc} [13] & [13] & [12] & [12] \\ [13] & [13] & [12] & [12] \\ [13] & [13] & [13] & [12] \\ [13] & [12] & [12] & [12] \end{array} \right] \end{array} \rightarrow \begin{array}{c} \left[\begin{array}{cccc} [13] & [13] & [12] & [12] \\ [13] & [12] & [12] & [12] \\ [13] & [13] & [13] & [12] \\ [13] & [12] & [12] & [12] \end{array} \right] \end{array} \rightarrow \dots \rightarrow \begin{array}{c} R \\ \left[\begin{array}{cccc} [1] & [1] & [1] & [1] \\ [2] & [1] & [1] & [0] \\ [1] & [2] & [2] & [0] \\ [1] & [4] & [1] & [0] \end{array} \right] \end{array}$$

Question: How to represent rows of the work matrix?

Linearization for lattice reduction of matrices

First consider Euclidean algorithm over $\mathbb{Z}/(7)[x]$

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Question: How to represent rows of the work matrix?

Idea: also used in earlier version of FLINT

$$\left[\begin{array}{cc|c} 3x^2 + 5x + 1 & 4x + 2 \\ 4x^2 + 2 & 6x + 2 \end{array} \right] \rightarrow \left[\begin{array}{ccc|ccc} 3 & 5 & 1 & 0 & 4 & 2 \\ 4 & 0 & 2 & 0 & 6 & 2 \end{array} \right]$$

From left equivalence to similarity

Recall companion matrix: 1×1 matrix of degree 4

$$[x^4 + 5x^3 + 6x^2 + 74x + 72] \longleftrightarrow xI_4 - \begin{bmatrix} & & & -72 \\ 1 & & & -74 \\ & 1 & & -6 \\ & & 1 & -5 \end{bmatrix}$$

Same idea works for matrices (sometimes): 2×2 matrix of degree 3

$$\left[\begin{array}{c|c} A & \\ \hline x^3 + 2x^2 + 6x + 6 & 4x^2 + 4x \\ 2x^2 + 5x + 3 & x^3 + 5x + 4 \end{array} \right] \longleftrightarrow xI_6 - \begin{array}{c} C \\ \left[\begin{array}{c|cc} & & -6 & 0 \\ & & -3 & -4 \\ \hline 1 & & -6 & -4 \\ & 1 & -5 & -5 \\ \hline & 1 & -2 & -4 \\ & & 1 & -2 & 0 \end{array} \right] \end{array}$$

Idea: Compute $\det A$ by computing Frobenius form of C in time $O((nd)^3)$

From left equivalence to similarity: general case

Input:

$$A = \begin{bmatrix} 5x^2 + 4x + 1 & x + 1 \\ 5x + 1 & 2x + 1 \end{bmatrix}$$

1. Random shift:

$$B = A \Big|_{x=x-2} = \begin{bmatrix} 5x^2 + 5x + 6 & x + 6 \\ 5x + 5 & 2x + 4 \end{bmatrix}$$

2. Revert:

$$C = x^2 B \Big|_{x=1/x} = \begin{bmatrix} 6x^2 + 5x + 5 & 6x^2 + x \\ 5x^2 + 5x & 4x^2 + 2x \end{bmatrix}$$

3. Normalize:

$$D = \begin{bmatrix} 6 & 6 \\ 5 & 4 \end{bmatrix}^{-1} \begin{bmatrix} 6x^2 + 5x + 5 & 6x^2 + x \\ 5x^2 + 5x & 4x^2 + 2x \end{bmatrix} = \begin{bmatrix} x^2 + 4x + 6 & 6x \\ 5x + 3 & x^2 \end{bmatrix}$$

Partial column linearization

$$\begin{array}{c}
 \begin{array}{c} A \\ \left[\begin{array}{ccc} [0] & & [5] \ [18] \\ & [0] & [5] \ [18] \\ & & [0] \ [5] \ [18] \\ & & & [6] \ [18] \\ & & & & [19] \end{array} \right]
 \end{array}
 \longleftrightarrow
 \begin{array}{c}
 \begin{array}{c} C \\ \left[\begin{array}{ccc|ccc} [0] & & [4] \ [4] & [0] & [4] & [4] \ [3] \\ & [0] & [4] \ [4] & [0] & [4] & [4] \ [3] \\ & & [0] \ [4] \ [4] & [0] & [4] & [4] \ [3] \\ & & & [1] & [4] & [4] \ [3] \\ & & & & [4] & [4] \ [4] \\ \hline & & -x^5 & 1 & & \\ & & & -x^5 & 1 & \\ & & & & -x^5 & 1 \\ & & & & & -x^5 & 1 \end{array} \right]
 \end{array}
 \end{array}
 \end{array}$$

- $\deg C =$ average column degree of A
- dimension of C is less than $2 \times$ dimension of A
- no computation required, only rewriting
- $\det C = \det A$, $\text{nullity}(C) = \text{nullity}(A)$, $C^{-1} = \left[\begin{array}{c|c} A^{-1} & * \\ \hline * & * \end{array} \right]$

Partial row and column linearization

$$A = \begin{bmatrix} [19] & [1] & [5] & [3] & [19] \\ [4] & [6] & [3] & [6] & [0] \\ [0] & [0] & [0] & [0] & [0] \\ [17] & [6] & [0] & [0] & [0] \\ [19] & [0] & [0] & [0] & [0] \end{bmatrix}$$

$$C = \left[\begin{array}{cccccccc|cccc} [4] & [1] & [4] & [3] & [4] & [4] & [4] & [4] & & -x^5 & & \\ [4] & [4] & [3] & [4] & [0] & & & & [1] & & & -x^5 \\ [0] & [0] & [0] & [0] & [0] & & & & & & & \\ [4] & [4] & [0] & [0] & [0] & [4] & [4] & [2] & [1] & & & \\ [4] & [0] & [0] & [0] & [0] & [4] & [4] & [4] & & & & \\ -x^5 & & & & & 1 & & & & & & \\ & & & & & -x^5 & 1 & & & & & \\ & & & & & & -x^5 & 1 & & & & \\ & -x^5 & & & & & & & 1 & & & \\ \hline & & [0] & & [4] & & & & & 1 & -x^5 & \\ & & & & [4] & & & & & & 1 & -x^5 \\ & & & & [4] & & & & & & & 1 \\ \hline & & & [1] & & & & & & & & 1 \end{array} \right].$$

- dimension of C is $< 3 \times (19 + 6 + 0 + 0 + 0)/5 = 15$